

Concept Parameters as a New Mechanism of Generic Programming for C# Language

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Generic Programming

A term “Generic Programming” (GP) was coined in 1989 by Alexander Stepanov and David Musser [1].

Idea

Code is written in terms of **abstract** types and operations (parametric polymorphism).

Purpose

Writing **highly reusable** code.

Constrained Generic Code

How to write a **generic** function that finds maximum element in a generic collection?

```
interface IEnumerable<T> : IEnumerable
{
    IEnumerator<T> GetEnumerator(); ... }

static T FindMax<T>(IEnumerable<T> vs) // could be ..(T[] vs)
{
    T mx = vs.First();
    foreach (var v in vs)
        if (mx < v)           // ERROR: operator '<'
            mx = v;         // is not provided for the type T
    ...
}
```

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```

To find maximum in vs, values of type **T** must be **comparable!**

“Being comparable” is a **constraint**.

An Example of Generic Code with Constraints (C#)

```
interface IEnumerable<T> : IEnumerable { ... }  
interface IComparable<T> { int CompareTo(T other); }  
  
static T FindMax<T>(IEnumerable<T> vs)  
    where T : IComparable<T> // F-bounded polymorphism  
{  
    T mx = vs.First();  
    foreach (var v in vs)  
        if (mx.CompareTo(v) < 0) mx = v;  
    return mx;  
}
```

Figure: Searching for maximum element in vs

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    return mx;
}
```

Figure: Searching for maximum element in vs

FindMax<T> can **only** be instantiated with types implementing the IComparable<T> interface.

```
var ints = new int[]{ 3, 2, -8, 61, 12 };
var iMax = FindMax(ints); // 61
var strs = new LinkedList<string>{ "hi", "bye", "stop", "hello" };
var sMax = FindMax(strs); // "stop"
```

Explicit Constraints on Type Parameters

Programming languages provide various language mechanisms for generic programming based on **explicit constraints**:

- Haskell: type classes;
- SML, OCaml: modules;
- Rust, Scala: traits;
- Swift: protocols;
- Ceylon, Kotlin, C#, Java: interfaces;
- etc.

C++

C++ Templates are unconstrained!

It was shown in earlier studies that C# and Java yield to many languages with respect to language support for GP [2–4].

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The Goal of the Research

To develop a mechanism of **generic programming** that improves language support for GP in mainstream **object-oriented** languages.

Motivation

Poor Language Support for Generic Programming

Is it a problem of C# and Java only?

Or is it a **typical** problem of **object-oriented** languages?

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What about **modern** object-oriented languages?

[name (first appeared, recent stable release)]

- Scala (2004, 2016);
- Rust (2010, 2016);
- Ceylon (2011, 2016);
- Kotlin (2011, 2016);
- Swift (2014, 2016).

Constraints-are-Types

All of them follow the same approach to constraining type parameters [5]: OO constructs used as **types** (such as *interfaces*) are also used as **constraints**.

Inevitable Limitations of the OO approach

(are usually “solved” with the Concept design pattern)

An interface/trait/protocol describes properties of a **single** type that implements/extends/adopts it. Therefore:

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Instead of

```
double Foo<A, B>(A[] xs) where <single constraint on A, B>  
// the constraint includes functions like B[] Bar(A a)
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Instead of

```
double Foo<A, B>(A[] xs) where <single constraint on A, B>
// the constraint includes functions like B[] Bar(A a)
```

we have:

```
interface IConstraintA<A, B> where A : IConstraintA<A, B>
where B : IConstraintB<A, B> {...}
interface IConstraintB<A, B> where A : IConstraintA<A, B>
where B : IConstraintB<A, B> {...}
double Foo<A, B>(A[] xs)
  where A : IConstraintA<A, B>
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double Foo<A, B>(A[] xs)
  where A : IConstraintA<A, B>
  where B : IConstraintB<A, B> {...}
```

- **Multiple models** cannot be supported at language level.

Related Work

There are several language extensions for generic programming influenced by Haskell type classes [6]:

- C++ concepts [7, 8] (2003–2014) and concepts in language G [9] (2005–2011);
- Generalized interfaces in JavaGI [10] (2007–2011);
- Constraints in Java Genus [11] (2015).

All these extensions follow the *alternative* approach to constraining type parameters.

The “Constraints-are-Not-Types” Approach

To **constrain** type parameters, a **separate** language construct is used. It cannot be used as type.

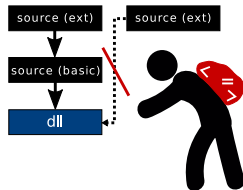
Drawbacks of the Existing Solutions

Neither of the extensions supports **all** the features:

- | | |
|------------------------|------------------------------|
| ① multiple models; | ④ supertype constraints; |
| ② associated types; | ⑤ concept-based overloading; |
| ③ subtype constraints; | ⑥ multiple dynamic dispatch. |

The extensions are implemented via **translation to the basic language**, but:

- ① resulting generic classes contain extra fields, whereas generic functions take extra arguments (this **brings run-time overhead**);
- ② translation is not reversible (this **breaks separate compilation**).



Research Track

- 1 To identify key **problems** of object-oriented languages with respect to their support for generic programming.
- 2 To **design** a **language extension** for C# that improves means of generic programming in the language.
- 3 To develop a **type-safe model** of the extension for FGJ [12].
- 4 To provide a “**proof-of-concept**” **implementation** of the extension for C#.

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Concepts and Generic Code in Cp#

(Cp# stands for C# extended with concept parameters)

Concepts:

```
concept Equality[T]
{
  bool Equal(T x, T y);
  bool NotEqual(T x, T y){ return !Equal(x, y); } }

```

```
concept Ordering[T] refines Equality[T]
{
  int Compare(T x, T y);
  bool Less(T x, T y) {...} ... }

```

```
concept Unifying[Tm, Eqtn, Subst]
{
  Subst Solve(IEnumerable<Eqtn> eqs); ... }

```

Generic Code:

```
bool Contains<T | Equality[T] eq>(IEnumerable<T> vs, T x)
{
  ... if (eq.Equal(...) )
}

```

```
interface ICollection<T> { ... bool Remove<|Equality[T] eq>(T x); }

```

```
class HashSet<T | Equality[T] eq> ...

```

Models in C#

```
// default case-sensitive equality comparison
model default EqStringCaseS for Equality[string] { ... }

// case-insensitive equality comparison
model EqStringCaseIS for Equality[string]
{
    bool Equal ( string x , string y )
    { return x.ToLower() == y.ToLower(); }
}

// default lexicographical ordering
model default OrdStringCSAsc for Ordering[string]
    refines EqStringCaseS { ... }
```

Models consistency is provided!

```
var s1 = new HashSet<string>(...); //s1 : HashSet<string|EqStringCaseS>
var s2 = new HashSet<string | EqStringCaseIS>(...);
s1.UnionWith(s2); // static ERROR: s1 and s2 have different types
```

Translation of Cp# to C#

- concept \Rightarrow generic interface
- model \Rightarrow class
- generic code \Rightarrow generic code with extra **type arguments**

```
static class ConceptSingleton<C> where C : new()
{ public static C Instance ... }
```

```
interface Equality<T> { bool Equal(T x, T y); }
interface Ordering<T> : Equality<T> { ... }
```

```
bool Contains<T, eq>(IEnumerable<T> vs, T x)
  where eq : Equality<T>, new()
{ ... if (ConceptSingleton<eq>.Instance.Equal(...) ...)
```

```
interface ISet<T, eq> where eq : Equality<T>, new() { ... }
class SortedSet<T, ord> : ISet<T, ord>
  where ord : Ordering<T>, new() { ... }
```

```
class EqStringCaseIS : Equality<string> { ... }
```

Benefits of the Translation Method

- 1 Extra **compile-time** type arguments are used instead of run-time class fields/function arguments.
- 2 Supplemented with attributes, the translation becomes **reversible!**

Why Is It Possible?

Because .NET CIL (Common Intermediate Language) preserves information on type parameters of generics.

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Comparison of Languages and Extensions

Language Support for GP in OO Languages	Haskell	C#	Java 8	Scala	Ceylon	Kotlin	Rust	Swift	JavaGI	G	C#cpt	Genus	ModImpl.
Constraints can be used as types	○	●	●	●	●	●	●	●	◐	○	○	○	○
<i>Explicit self types</i>	—	○	○	◐	●	○	●	●	◐	—	—	—	—
Multi-type constraints	●	*	*	*	○	*	○	○	●	●	●	●	●
<i>Retroactive type extension</i>	—	●	○	○	○	●	●	●	○	○	○	○	—
<i>Retroactive modeling</i>	●	*	*	*	○	*	●	●	●	●	●	●	●
<i>Type conditional models</i>	●	○	○	○	○	○	●	○	●	●	●	●	●
<i>Static methods</i>	●	○	●	○	●	●	●	●	●	●	●	●	●
<i>Default method implementation</i>	●	○	●	●	●	●	●	●	◐	●	●	○	○
<i>Associated types</i>	●	○	○	●	○	○	●	●	○	●	●	○	●
<i>Constraints on associated types</i>	◐	—	—	●	—	—	●	●	—	●	●	—	●
<i>Same-type constraints</i>	◐	—	—	●	—	—	●	●	—	●	●	—	●
<i>Concept-based overloading</i>	○	○	○	○	○	○	●	○	○	◐	○	○	○
Multiple models	○	*	*	*	*	*	○	○	○	◐ ^a	●	●	●
Models consistency (model-dependent types)	— ^b	○	○	○	○	○	— ^b	— ^b	— ^b	— ^b	●	●	●
<i>Model genericity</i>	—	*	*	*	*	*	●	○	○	○	○	●	—
<i>Multiple dynamic dispatch</i>	—	○	○	○	○	○	○	○	◐	○	○	●	—

* means support via the Concept pattern. ^aG supports lexically-scoped models but not really multiple models.

^bIf multiple models are not supported, the notion of model-dependent types does not make sense.

Concept Pattern I

With the Concept design pattern [13] (“Type Classes As Objects and Implicits” by Oliveira et. al., 2010), constraints on type parameters are replaced with extra arguments – “**concepts**”.

F-Bounded Polymorphism

```

interface IComparable<T>
{ int CompareTo(T other); } // *

static T FindMax<T>(
    IEnumerable<T> vs)
where T : IComparable<T> // *
{
    T mx = vs.First();
    foreach (var v in vs)
        if (mx.CompareTo(v) < 0) // *
            ...

```

Concept Pattern

```

interface IComparer<T>
{ int Compare(T x, T y); } // *

static T FindMax<T>(
    IEnumerable<T> vs,
    IComparer<T> cmp) // *
{
    T mx = vs.First();
    foreach (var v in vs)
        if (cmp.Compare(mx,v) < 0) // *
            ...

```

Concept Pattern II

In Scala it has a special support: **context bounds** and **implicit**s.

F-Bounded Polymorphism

```

trait Ordered[A] {
  abstract def compare
    (that: A): Int
  def < (that: A): Boolean = ...
}

// upper bound
def findMax[A <: Ordered[A]]
  (vs: Iterable[A]): A
{ ... }

```

Concept Pattern

```

trait Ordering[A] {
  abstract def compare
    (x: A, y: A): Int
  def lt(x: A, y: A): Boolean = ...
}

// context bound (syntactic sugar)
def findMax[A : Ordering]
  (vs: Iterable[A]): A
{ ... }

// implicit argument (real code)
def findMax(vs: Iterable[A])
  (implicit ord: Ordering[A])
{ ... }

```


Advantages of the Concept Pattern

Both limitations of the “Constraints-are-Types” approach are eliminated with this design pattern!

- 1 multi-type constraints are multi-type “concept” arguments;

```
interface IConstraintAB<A, B>
{ B[] Bar(A a); ... }
```

```
double Foo<A, B>(A[] xs, IConstraintAB<A, B> c)
{ ... c.Bar(...) ... }
```

- 2 multiple “models” are allowed as long as several classes can implement the same interface.

```
class IntCmpDesc : IComparer<int> { ... }
class IntCmpMod42 : IComparer<int> { ... }
```

```
var ints = new int[]{ 3, 2, -8, 61, 12 };
```

```
var minInt = FindMax(ints, new IntCmpDesc());
var maxMod42 = FindMax(ints, new IntCmpMod42());
```

Drawbacks of the Concept Pattern

The Concept design pattern is widely used in standard generic libraries of C#, Java, and Scala, but it has serious problems.

Drawbacks

- 1 runtime overhead (extra class fields or function arguments);
- 2 models inconsistency.

```
interface IEqualityComparer<T>
{ ... }
```

```
class HashSet<T> : ...
{
    IEqualityComparer<T>
        Comparer;
    ...
}
```

```
static HashSet<T> GetUnion<T>
    (HashSet<T> a, HashSet<T> b)
{ var us = new HashSet<T>
    (a, a.Comparer);
  us.UnionWith(b);
  return us; }
```

Attention! GetUnion(s1, s2)
could differ from
GetUnion(s2, s1)!